

# (Quasi) Newton methods

## Theory and sparse implementation

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# Summary

- Nonlinear systems of equations. A few examples
- Newton's method for  $f(x) = 0$ .
- Newton's method for systems.
- Local convergence. Exit tests.
- Global convergence. Backtracking. Line search algorithms
- Two computationally useful variants: the Inexact Newton method and the Quasi Newton method.

# Sistemi di equazioni non lineari

## Some examples

- Intersection of curves in  $\mathbb{R}^n$ .

Find the intersections between the circumference and the hyperbola:

$$\begin{cases} x^2 + y^2 = 4 \\ xy = 1 \end{cases}$$

- Equation describing the flow in porous media (Richards' equation):

$$\frac{\partial \psi}{\partial t} - \vec{\nabla} \cdot \left( K(\psi) \vec{\nabla} \psi \right) = f \quad (1)$$

- Function minimization (applications in data science, machine learning)

$$\min G(x) \implies \text{Solve } G'(x) = 0$$

## Newton's method

Given a function  $f \in C^1$ , we aim at finding one solution of the equation

$$f(x) = 0$$

Given  $x_k$ , an approximation to the solution  $\xi$ , we correct it to find  $x_{k+1} = x_k + s$

We impose the condition  $f(x_{k+1}) = 0$  and expand  $f(x_{k+1})$  in Taylor series neglecting the terms of order greater or equal than 2.

$$0 = f(x_{k+1}) = f(x_k) + sf'(x_k)$$

from which

$$s = -\frac{f(x_k)}{f'(x_k)}.$$

The Newton's method can therefore be written as

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$

# Newton's method for system of nonlinear equations

Let us now solve the following nonlinear system

$$\begin{cases} F_1(x_1, x_2, \dots, x_n) &= 0 \\ F_2(x_1, x_2, \dots, x_n) &= 0 \\ \dots &= 0 \\ F_n(x_1, x_2, \dots, x_n) &= 0 \end{cases} \quad (2)$$

more synthetically

$$F(x) = 0$$

where

$$F = \begin{pmatrix} F_1 \\ F_2 \\ \dots \\ F_n \end{pmatrix} \quad x = \begin{pmatrix} x_1 \\ x_2 \\ \dots \\ x_n \end{pmatrix}$$

Let us assume that  $F$  be differentiable in an open subset of  $\mathbb{R}^n$ .

## Newton's method for system of nonlinear equations

As in the scalar case, we try to correct an approximation  $x_k$  as  $x_{k+1} = x_k + s$ .

Let us impose  $F(x_{k+1}) = 0$  and as before expand in Taylor series the function  $F(x_{k+1})$ .

$$0 = F(x_{k+1}) = F(x_k) + F'(x_k)s$$

where  $F'(x_k)$  is the Jacobian of system (7) evaluated in  $x_k$  i. e.

$$(F'(x))_{ij} = \frac{\partial F_i}{\partial x_j}(x)$$

As before the problem is to compute the increment  $s$  which is now a vector of  $n$  components.

$$s = -(F'(x_k))^{-1}F(x_k)$$

The  $k$ th iteration of the Newton's method is thus written as

$$x_{k+1} = x_k - (F'(x_k))^{-1}F(x_k)$$

# Newton's method for system of nonlinear equations

Some observations:

- The Jacobian matrix  $F'(x_k)$  must be invertible.
- Local convergence of the Newton's method can be proved provided that the initial approximation  $x_0$  is sufficiently close to the solution.
- Computation of  $x_{k+1}$  starting from  $x_k$  requires inversion of (possibly large and sparse) Jacobian matrix. This operation is inefficient as known. In practice vector  $s$  is evaluated by solving the following linear system

$$F'(x_k)s = -F(x_k)$$

- $F'$  is often non symmetric, so GMRES iterative method is suggested for the solution of the Newton system

# Algorithm

Let us write a first version of the Algorithm, by taking into account previous comments.

## Newton Algorithm

Given an initial approximation  $x_0$  ,  $k := 0$ .

repeat until convergence

- solve:  $F'(x_k)s = -F(x_k)$
- $x_{k+1} := x_k + s$
- $k := k + 1$



# Newton method:convergence

## Standard Assumptions.

- Equation (1) has a unique solution which we denote with  $x^*$ .
- The Jacobian  $F'$  is Lipschitz continuous. There exists a real scalar  $\gamma$  such that

$$\|F'(y) - F'(x)\| \leq \gamma \|y - x\|$$

- $F'(x^*)$  is nonsingular.

**Notations.** Let us define:

- the error at the iteration  $k$ :  $e_k = x_k - x^*$

## Theorem

*There exists  $\delta > 0$  such that*

$$\|e_0\| < \delta \implies \|e_{k+1}\| \leq K \|e_k\|^2$$

*with*

$$K = \gamma \|(F'(x^*))^{-1}\|$$

## Proof of Newton convergence

We premise the following

### Theorem

Let  $F$  be differentiable in an open set  $\Omega \in \mathbb{R}^n$ , and  $\mathbf{x}^* \in \Omega$ . Then for  $\mathbf{x} \in \Omega$  sufficiently close to  $\mathbf{x}^*$

$$F(\mathbf{x}) - F(\mathbf{x}^*) = \int_0^1 F'((\mathbf{x}^* + t(\mathbf{x} - \mathbf{x}^*)))(\mathbf{x} - \mathbf{x}^*) dt$$

### Proof.

Let  $g(t) = F((\mathbf{x}^* + t(\mathbf{x} - \mathbf{x}^*)))$ . Using the chain rule

$$g'(t) = F'((\mathbf{x}^* + t(\mathbf{x} - \mathbf{x}^*)))(\mathbf{x} - \mathbf{x}^*)$$

Hence by the Fundamental Theorem of Calculus

$$g(1) - g(0) = \int_0^1 g'(t) dt = F(\mathbf{x}) - F(\mathbf{x}^*).$$



## Proof of Newton convergence

We now need two further results. The first one is also known as *Banach Lemma*

### Lemma

Let  $A, B$  square  $n \times n$  matrices, and  $B$  such that  $\|I - BA\| < 1$ . Then  $A, B$  are both nonsingular and

$$\|A^{-1}\| \leq \frac{\|B\|}{1 - \|I - BA\|}.$$

### Lemma

Let the standard assumptions hold. Then there is  $\delta > 0$  so that for all  $x \in \mathcal{B}(\delta) = \{x : \|x - x^*\| < \delta\}$  the following hold true:

$$\|F'(x)\| \leq 2\|F'(x^*)\| \quad (3)$$

$$\|F'(x)^{-1}\| \leq 2\|F'(x^*)^{-1}\| \quad (4)$$

## Proof of Newton convergence

Proof.

(3). By triangular inequality and Lipschitz continuity

$$\|F'(x)\| - \|F'(x^*)\| \leq \|F'(x) - F'(x^*)\| \leq \gamma \|x - x^*\| = \gamma \|e\| \leq \gamma \delta.$$

Now if  $\delta < \frac{\|F'(x^*)\|}{\gamma}$  then  $\|F'(x)\| \leq \|F'(x^*)\| + \gamma \delta \leq 2\|F'(x^*)\|$ .

(4). Choosing now  $\delta < \frac{1}{2\gamma\|F'(x^*)^{-1}\|}$  then

$$\|I - F'(x^*)^{-1}F'(x)\| = \|F'(x^*)^{-1}\| \|F'(x^*) - F'(x)\| \leq \|F'(x^*)^{-1}\| \gamma \|e\| \leq \frac{1}{2}.$$

Then applying Banach's Lemma

$$\|F'(x)^{-1}\| \leq \frac{\|F'(x^*)^{-1}\|}{1 - \|I - F'(x^*)^{-1}F'(x)\|} \leq \frac{\|F'(x^*)^{-1}\|}{1 - 1/2} = 2\|F'(x^*)^{-1}\|.$$



## Proof of Newton convergence

We are now able to prove the main theorem.

Proof.

Let  $\delta$  be smaller enough so that the previous Lemma holds.

$$\begin{aligned}e_{k+1} &= x_{k+1} - x^* = x_k - x^* - F'(x_k)^{-1} F(x_k) \\&= e_k - F'(x_k)^{-1} \int_0^1 F'((x^* + t(x_k - x^*))) e_k dt \\&= F'(x_k)^{-1} \int_0^1 (F'(x_k) - F'((x^* + t(x_k - x^*)))) e_k dt.\end{aligned}$$

Taking norms and again using Lipschitz continuity

$$\begin{aligned}\|e_{k+1}\| &\leq \|F'(x_k)^{-1}\| \int_0^1 \gamma \|x_k - x^* - t(x_k - x^*)\| \|e_k\| dt \\&= \|F'(x_k)^{-1}\| \gamma \|e_k\|^2 \int_0^1 (1-t) dt = \gamma \|F'(x^*)^{-1}\| \|e_k\|^2 = K \|e_k\|^2.\end{aligned}$$



## Exit test

Ideal exit test  $\|e_k\| < \varepsilon$  (absolute error)

or  $\|e_k\| < \varepsilon \|e_0\|$  (relative error); where  $\varepsilon$  is a prescribed tolerance.

As however  $x^*$  is not known

- Exit test on the relative residual. Stop when

$$\frac{\|F(x_k)\|}{\|F(x_0)\|} < \varepsilon$$

- Exit test on the difference. Stop when

$$\|s\| = \|x_{k+1} - x_k\| < \varepsilon$$

## Exit tests

### Motivations

- Test on the residual. It can be proved that for a sufficiently small  $\delta$

$$\frac{1}{4\kappa} \frac{\|e_k\|}{\|e_0\|} \leq \frac{\|F(x_k)\|}{\|F(x_0)\|} \leq 4\kappa \frac{\|e_k\|}{\|e_0\|}$$

where  $\kappa = \|F'(x^*)\| \|(F'(x^*))^{-1}\|$  is the condition number of  $F'(x^*)$ .

If  $F'(x^*)$  is well-conditioned ( $\kappa \approx 1$ ), the test is reliable.

- Test on the difference

$$\begin{aligned} x_{k+1} - x_k &= x_{k+1} - x^* + x^* - x_k = e_{k+1} - e_k \\ \|x_{k+1} - x_k\| &= \|e_k\| + O(\|e_k\|^2) \end{aligned}$$

The difference at step  $k+1$  is of the same order of magnitude as the error at step  $k$ .  
(Pessimistic test)

## Example

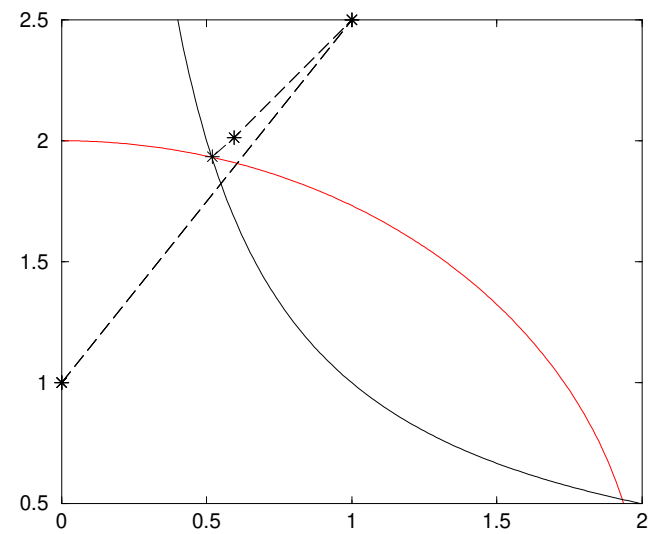
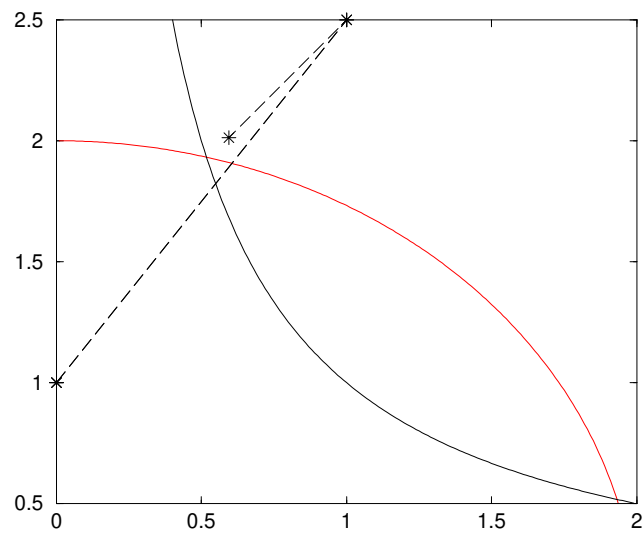
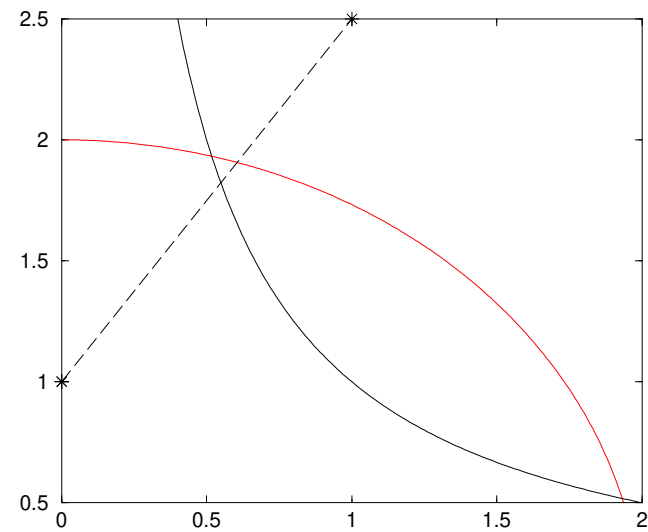
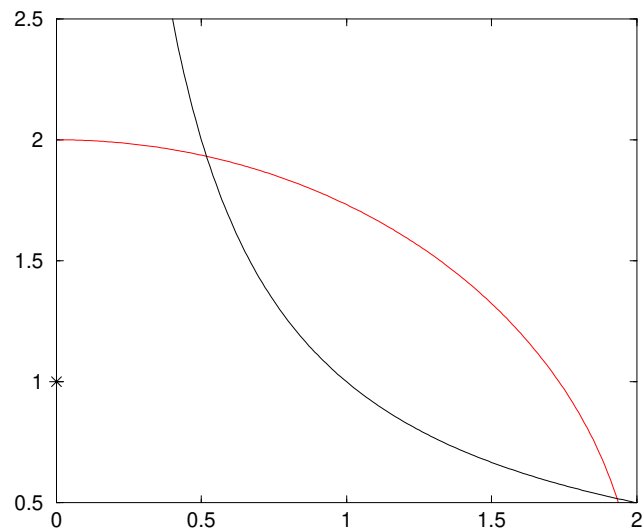
$$\begin{cases} x^2 + y^2 - 4 = 0 \\ xy - 1 = 0 \end{cases} \quad x^{(0)} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$k$	$x_1^{(k)}$	$x_2^{(k)}$	$\ e^{(k)}\ $	$\ x^{(k+1)} - x^{(k)}\ $	$\frac{\ e^{(k+1)}\ }{\ e^{(k)}\ ^2}$
0	0.000000000	1.000000000	$0.107 \times 10^{+01}$		
1	1.000000000	2.500000000	$0.745 \times 10^{+00}$	$0.180 \times 10^{+01}$	0.655899
2	0.595238095	2.011904761	$0.111 \times 10^{+00}$	$0.634 \times 10^{+00}$	0.200716
3	0.520020336	1.934236023	$0.337 \times 10^{-02}$	$0.108 \times 10^{-00}$	0.271153
4	0.517640404	1.931853966	$0.327 \times 10^{-05}$	$0.337 \times 10^{-02}$	0.288114
5	0.517638090	1.931851652	$0.309 \times 10^{-11}$	$0.327 \times 10^{-05}$	0.288656

$$\|F(x^{(0)})\| = 3.16 \quad \|F(x^{(1)})\| = 3.58$$



# Example



# Global Convergence

- Convergence of Newton's method not guaranteed. Frequently Newton's step moves away from the solution
- To avoid divergence we accept Newton's step if the following condition holds:  
 $\|F(x_{k+1})\| < \|F(x_k)\|$
- If the above condition is not satisfied, then the Newton step is reduced  $\implies$  "backtracking" or "linesearch".

## Algorithm: Newton 2.

Given an initial approximation  $x_0$ ,  $k := 0$ .

repeat until convergence

- solve:  $F'(x_k)s = -F(x_k)$
- $x_t := x_k + s$
- if  $\|F(x_t)\| < \|F(x_k)\|$  then  $x_{k+1} := x_t$
- else  $s := s/2$ , go to (•)
- $k := k + 1$

# Esempio

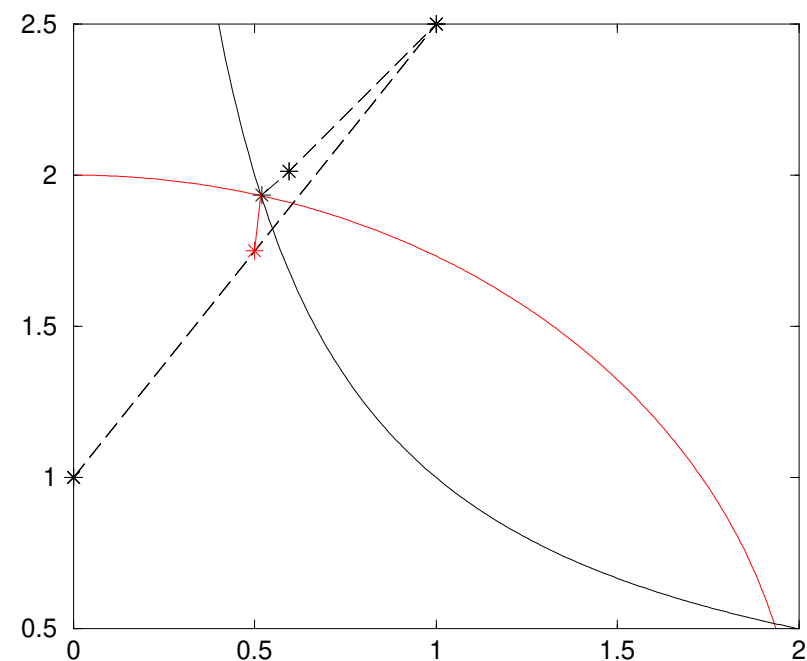
## Newton con backtracking

$$\begin{cases} x^2 + y^2 - 4 = 0 \\ xy - 1 = 0 \end{cases} \quad x^{(0)} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$k$	$x_1^{(k)}$	$x_2^{(k)}$	$\ e^{(k)}\ $	$x^{(k+1)} - x^{(k)}$	$\frac{\ e^{(k+1)}\ }{\ e^{(k)}\ ^2}$
0	0.000000000	1.000000000	$0.106597 \times 10^{+01}$		
1	0.500000000	1.750000000	$0.182705 \times 10^{+00}$	0.901388E+01	0.160790

$$\|F(x^{(0)})\| = 3.16$$

$$\|F(x^{(1)})\| = 0.699$$



## Inexact Newton Methods (Hint)

- Idea: Avoid **Oversolving** linear systems when too far from the nonlinear solution.
- **Example.** FE Steady-state Richards Equation

$$A(\psi)\psi = b(\psi), \quad F(x) = A(x)x - b(x), \quad F' = A + \frac{\partial(A)}{\partial x}x$$

- $F'$  inherit the same size and sparsity of  $A$ .
- Generic Newton iteration:
  - Solve:  $F'(\mathbf{x}_k)\mathbf{s} = -F(\mathbf{x}_k)$
  - $\mathbf{x}_{k+1} := \mathbf{x}_k + \mathbf{s}$
  - $k := k + 1$
- We want to solve the linear system by an iterative method with **variable tolerance**.

$$\|F'(\mathbf{x}_k)\mathbf{s} + F(\mathbf{x}_k)\| \leq \eta_k \|F(\mathbf{x}_k)\|$$

- $\eta_k$  can be “large” at the beginning of Newton iteration and must be small towards the end.
- If  $\eta_k \rightarrow 0$  then Newton convergence is **superlinear**.
- If  $\eta_k = O(\|F(\mathbf{x}_k)\|)$  then again quadratic convergence can be proved.

# Quasi-Newton Methods

Motivation: Jacobian matrix

- Not always explicitly available (sometimes function  $F$  is known as a set of data)  
or
- Differentiation of  $F$  may be too costly to be afforded at every Newton iteration

A possible answer to this problem is given by the quasi-Newton methods which compute a sequence of approximate Jacobians possibly starting from the 'true' initial Jacobian.

Instead of solving

$$x_{k+1} = x_k - F'(x_k)^{-1} F(x_k)$$

we solve

$$x_{k+1} = x_k - B_k^{-1} F(x_k)$$

## Quasi-Newton Methods

Sequence of  $B_k$  can be constructed in many ways. The simplest approach is due to Broyden:

$$B_{k+1} = B_k + \frac{(y - B_k s)s^T}{s^T s} = B_k + \frac{F(\mathbf{x}_{k+1})s^T}{s^T s}$$

where  $y = F(\mathbf{x}_{k+1}) - F(\mathbf{x}_k)$  and using  $B_k s_k = -F(\mathbf{x}_k)$ .

The Broyden update formula satisfies:

- ① the secant condition, namely  $B_{k+1}s = y$ .
- ②  $B_{k+1}$  is the closest matrix to  $B_k$  in the Frobenius norm among all the matrices satisfying the secant condition.

$$B_{k+1} = \underset{B : Bs = y}{\operatorname{argmin}} \|B - B_k\|$$

**NOTE:** Frobenius norm of a matrix is defined as

$$\|A\|_F = \sqrt{\sum_i \sum_j a_{ij}^2} = \operatorname{tr}(A^T A)$$

## Quasi-Newton Methods

The secant condition  $B_{k+1}s = y$  is a generalization of the secant method (Regula falsi):

$$x_{k+1} = x_k - \frac{f(x_k)}{b_k}$$

where

$$b_k = \frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}} = \frac{y}{s}.$$

- In  $n$  dimension, the secant condition has an infinite number of solutions.
- Broyden's choice satisfies secant condition, in fact

$$B_{k+1}s = B_k s + \frac{(y - B_k s)s^T}{s^T s} s = B_k s + y - B_k s = y.$$

- Broyden's choice satisfies the **least change condition**.

$$\begin{aligned} \|B_{k+1} - B_k\| &= \left\| \frac{(y - B_k s)s^T}{s^T s} \right\|_F = \left\| \frac{(B s - B_k s)s^T}{s^T s} \right\|_F = \\ &= \left\| \frac{(B - B_k) s s^T}{s^T s} \right\|_F \leq \|B - B_k\|_F \left\| \frac{s s^T}{s^T s} \right\|_F = \|B - B_k\|_F. \end{aligned}$$

using  $\left\| \frac{s s^T}{s^T s} \right\|_F = 1$  (proof by exercise).

# Convergence Results

## Definition

A sequence  $\{x_k\}$  converges superlinearly to  $x^*$  if there are  $\alpha > 1$  and  $K > 0$  such that

$$\|x_{k+1} - x^*\| \leq K \|x_k - x^*\|^\alpha$$

Let us now define the error in jacobian approximations:

$$E_k = B_k - F'(x^*)$$

The first Theorem states that the difference between the exact and the approximate Jacobian does not grow with the Newton iteration. This property is also called *bounded deterioration*.

## Theorem

$$\|E_{k+1}\| \leq \|E_k\| + \frac{\gamma}{2} (\|e_k\| + \|e_{k+1}\|)$$



# Convergence Results and implementation

## Theorem

*Let the standard assumption holds. Then there are  $\delta$  and  $\delta_B$  such that if  $\|e_0\| < \delta$  and  $\|E_0\| < \delta_B$  the Broyden sequence exists and  $x_n \rightarrow x^*$  superlinearly.*

This theorem states that we can make  $\|E_k\|$  as small as we want by properly choosing the initial vector  $x_0$  and the initial Jacobian approximation  $B_0$ .

If it is the case, the convergence of the iteration remains very fast (superlinear convergence).

## Problem.

How to implement solution of Newton system with  $B_k^{-1}$  instead of  $J(x_k)^{-1}$  ?

Note that even if  $B_0$  is sparse  $B_1$  is not.

Careful implementation should avoid inversion of dense matrices.

## Sparse implementation of Broyden method: Sherman Morrison formula.

### Theorem

$$(B + uv^T)^{-1} = \left( I - \frac{(B^{-1}u)v^T}{1 + v^T B^{-1}u} \right) B^{-1} = B^{-1} - \frac{(B^{-1}u)v^T B^{-1}}{1 + v^T B^{-1}u}$$

### Proof.

Let us look for the inverse of  $B + uv^T$  as  $B^{-1} + xy^T$ . The following conditions must hold:

$$(1) \quad I = (B + uv^T) \cdot (B^{-1} + xy^T) = I + uv^T B^{-1} + Bxy^T + uv^T xy^T$$

$$(2) \quad I = (B^{-1} + xy^T) \cdot (B + uv^T) = I + xy^T B + B^{-1}uv^T + xy^T uv^T$$

Multiplying the first by  $u$  on the right and the second by  $v^T$  on the left yields:

$$u \left( v^T B^{-1} u \right) + Bx(y^T u) + u(v^T xy^T u) = 0 \quad \implies \quad x = \alpha B^{-1} u.$$

$$(v^T x)y^T B + (v^T B^{-1} u)v^T + (v^T xy^T u)v^T = 0 \quad \implies \quad y = \beta B^{-1} v.$$

Without loss of generality set  $\beta = 1$  hence substituting in (1) we obtain

$$0 = uv^T B^{-1} + \alpha uv^T B^{-1} + uv^T \alpha B^{-1} uv^T B^{-1} = uv^T B^{-1} \left( 1 + \alpha(1 + v^T B^{-1} u) \right)$$

Finally we get  $\alpha = \frac{-1}{1 + v^T B^{-1} u}$  which completes the proof. □

## Sparse implementation of Broyden method

Need to compute  $B_k^{-1}F(x_k)$  without

- 1 Computing  $B_k^{-1}$  since we do not want to invert matrices.
- 2 Computing  $B_k$  since it is dense.

In our context we need to evaluate  $B_{k+1}^{-1}$  in terms of  $B_k^{-1}$  starting from

$$B_{k+1} = B_k + u_k v_k,$$

where we can define among the others

$$u_k = \frac{F(x_{k+1})}{\|s_k\|}, \quad v_k = \frac{s_k}{\|s_k\|}, \quad \text{so that}$$

## Sparse implementation of Broyden method

$$\begin{aligned} B_{k+1}^{-1} &= (B_k + u_k v_k^T)^{-1} = \left( I - \frac{(B_k^{-1} u_k) v_k^T}{1 + v_k^T B_k^{-1} u_k} \right) B_k^{-1} \\ &= \left( I - w_k v_k^T \right) B_k^{-1} \end{aligned}$$

Where we have defined  $w_k = \frac{B_k^{-1} u_k}{1 + v_k^T B_k^{-1} u_k}$ . Now by induction

$$B_k^{-1} = \left( I - w_{k-1} v_{k-1}^T \right) \left( I - w_{k-2} v_{k-2}^T \right) \cdots \left( I - w_0 v_0^T \right) B_0^{-1}$$

# Sparse implementation of Broyden method

Important results:  $s_k = -B_k^{-1}F_k$  is accomplished by

- ① Solving the system  $B_0 z_0 = -F_k$
- ② Computing  $\alpha_0 = w_0^T z_0$ , then  $z_1 = z_0 - \alpha_0 w_0$   
Computing  $\alpha_1 = w_1^T z_1$ , then  $z_2 = z_1 - \alpha_1 w_1$   
...  
Computing  $\alpha_{k-1} = w_{k-1}^T z_{k-1}$ , then  $z_k = z_{k-1} - \alpha_{k-1} w_{k-1}$

**Problem.** We do not know how to compute  $w_j, j = 1, \dots, k-1$ . Let us define

$$p = B_{k-1}^{-1}F(x_k) = (I - w_{k-2}v_{k-2}^T) \cdots (I - w_0v_0^T) B_0^{-1}F(x_k)$$

It follows that

$$\begin{aligned} s_k &= -B_k^{-1}F_k = -\left(I - w_{k-1}v_{k-1}^T\right)p = w_{k-1}(v_{k-1}^T p) - p \\ B_{k-1}^{-1}u_{k-1} &= B_{k-1}^{-1} \frac{F_k}{\|s_{k-1}\|} = \frac{p}{\|s_{k-1}\|} \\ w_{k-1} &= \frac{B_{k-1}^{-1}u_{k-1}}{1 + v_{k-1}^T B_{k-1}^{-1}u_{k-1}} = \frac{p}{\|s_{k-1}\| + v_{k-1}^T p} \end{aligned}$$

## Sparse implementation of Broyden method

Now combining  $s_k = w_{k-1}(v_{k-1}^T p) - p$  with  $w_{k-1} = \frac{p}{\|s_{k-1}\| + v_{k-1}^T p}$  we obtain

$$\|s_{k-1}\| w_{k-1} = p - w_{k-1} v_{k-1}^T p = -s_k$$

hence

$$w_{k-1} = -\frac{s_k}{\|s_{k-1}\|}$$

Hence  $B_k^{-1}$  can be written in terms of sequence  $\{s_j\}$  only as

$$B_k^{-1} = \prod_{j=0}^{k-1} (I - w_j v_j) = \prod_{j=0}^{k-1} \left( I + \frac{s_{j+1} s_j^T}{\|s_j\|_2^2} \right)$$

**NOTE:** We know  $s_k$  as a function of  $B_k^{-1}$  and  $B_k^{-1}$  as a function of  $s_k$ .

## Sparse implementation of Broyden method

Let us write  $s_{k+1}$  as

$$\begin{aligned} s_{k+1} = -B_{k+1}^{-1} F_{k+1} &= -\left(I + \frac{s_{k+1} s_k^T}{\|s_k\|_2^2}\right) \prod_{j=1}^{k-1} \left(I + \frac{s_{j+1} s_j^T}{\|s_j\|_2^2}\right) B_0^{-1} F_k \\ &= -\left(I + \frac{s_{k+1} s_k^T}{\|s_k\|_2^2}\right) B_k^{-1} F_{k+1} \end{aligned}$$

or

$$s_{k+1} = -B_k^{-1} F_{k+1} - s_{k+1} \frac{s_k^T B_k^{-1} F_{k+1}}{\|s_k\|_2^2}$$

Finally solve for  $s_{k+1}$  to obtain

$$s_{k+1} = \frac{-B_k^{-1} F_{k+1}}{1 + s_k^T B_k^{-1} F_{k+1} / \|s_k\|_2^2}$$

## Broyden Algorithm (sketch)

- **INPUT**:  $x_0, B_0$ . Set  $k := 0, x := x_0$ .
- First step: **Solve**  $B_0 s_0 = -F(x_0)$
- **repeat** until convergence
  - $x := x + s_k$
  - **Solve**  $B_0 z = -F(x)$
  - $k := k + 1$ .
  - **for**  $j := 0$  **to**  $k - 1$ 
$$z := z + s_{j+1} \frac{s_j^T z}{\|s_j\|_2^2}$$
  - **end for**
  - $s_{k+1} := \frac{z}{1 - s_k^T z / \|s_k\|_2^2}$
- **end repeat**

- 1 Only a system with  $B_0$  needed to be solved at each Newton step.
- 2  $k$  scalar products and 1 vector norm at  $k$ -th step. Complexity increasing with  $k$ .